splitMEM: graphical pan-genome analysis with suffix skips

Shoshana Marcus May 7, 2014

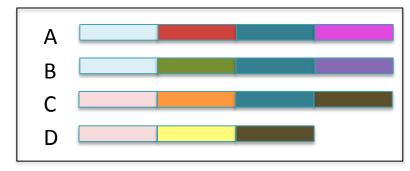


Outline

- I Overview
- 2 Data Structures
- 3 splitMEM Algorithm
- 4 Pan-genome Analysis

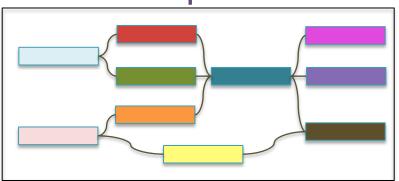
Objective

Input



- Several complete genomes
- Available today for many microbial species, near future for higher eukaryotes
- Pan-genome: analyze multiple genomes of species together

Output

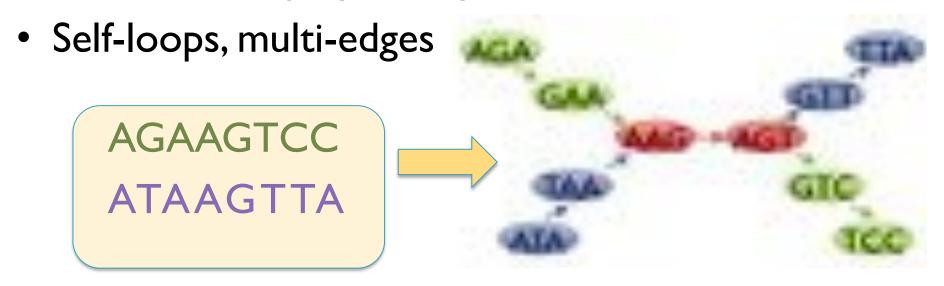


Compressed de Bruijn graph

- Graphical representation depicts how population variants relate to each other, especially where they diverge at branch points
- How well conserved is a sequence?
- What are network properties?

de Bruijn graph

- Node for each distinct kmer
- Directed edge connects consecutive kmers
- Nodes overlap by k-1 bp



Reconstruct original sequence:

Eulerian path through graph, visit each edge once

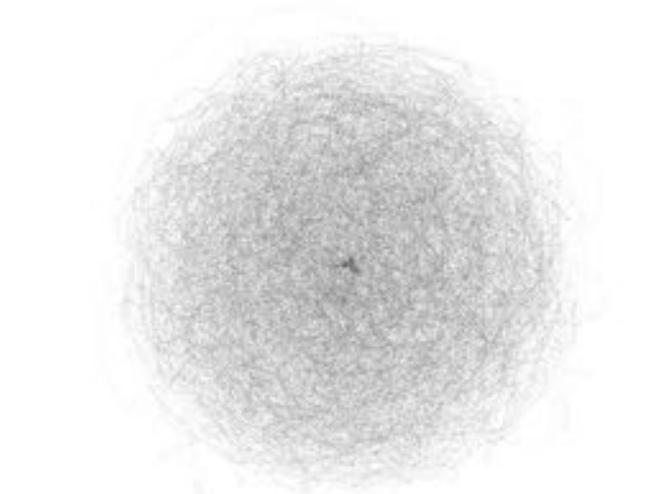
Compressed de Bruijn graph

- Merge non-branching chains of nodes
- Min. number of nodes that preserve path labels



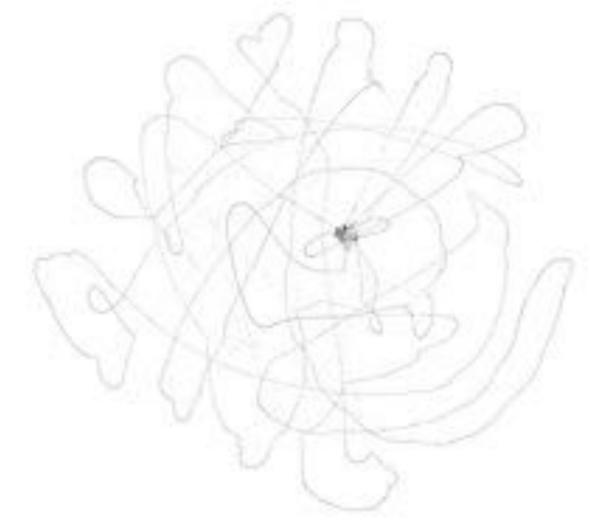
- ♦Usually built from uncompressed graph
- ♦ We build directly in O(n log n) time and space

Compresssed de Bruijn graph



9 strains of Bacillus anthracis k=25

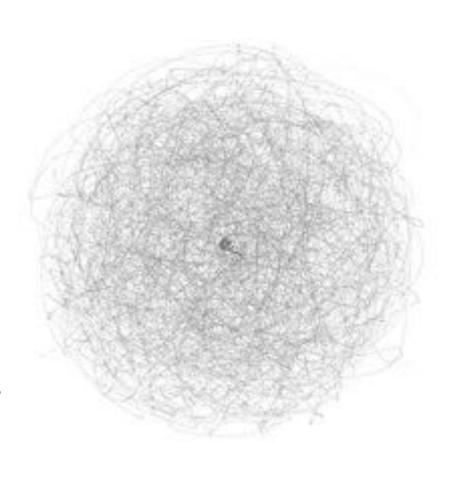
Compresssed de Bruijn graph



9 strains of Bacillus anthracis k=1000

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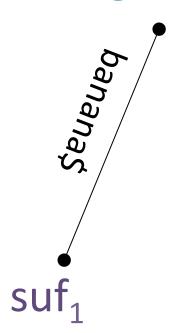
Suffix Tree

- Rooted, directed tree with leaf for each suffix.
- Each internal node, except the root, has at least two children.
- Each edge is labeled with nonempty substring.
- No two siblings begin with the same character.
- Path from root to leaf i spells suffix S[i ... n].
- Append special character \$ to guarantee each suffix ends at leaf.





Naïve Algorithm

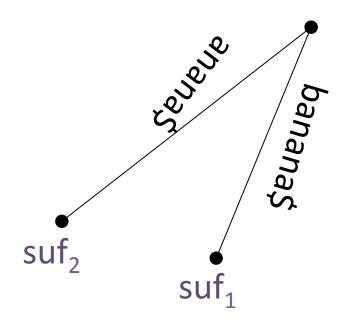


S = banana\$

suf₁ banana\$



Naïve Algorithm

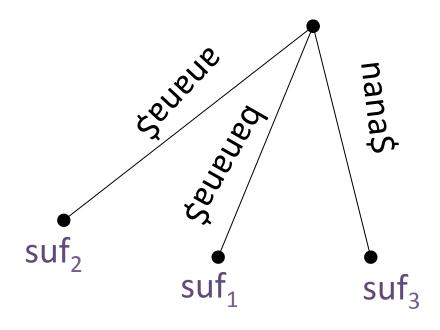


S = banana\$

suf₂ anana\$



Naïve Algorithm

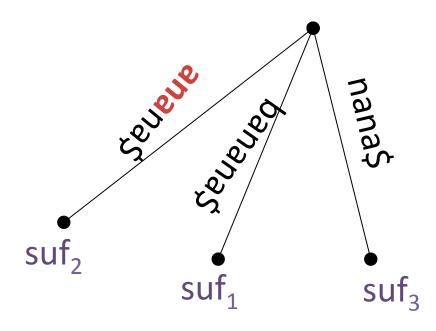


S = banana\$

suf₃ nana\$



Naïve Algorithm

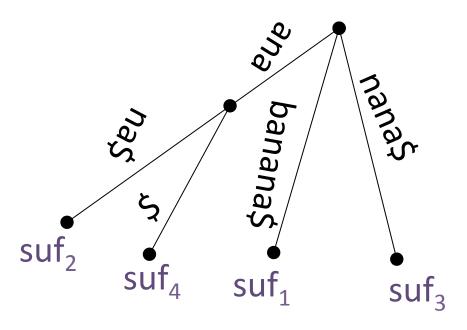


S = banana\$

suf₄ ana\$



Naïve Algorithm

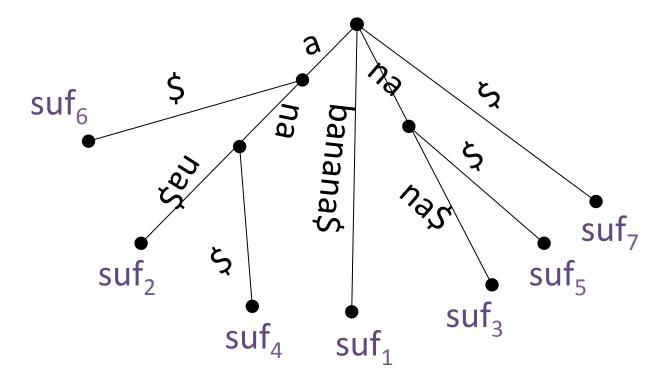


S = banana\$

suf₄ ana\$



Naïve Algorithm

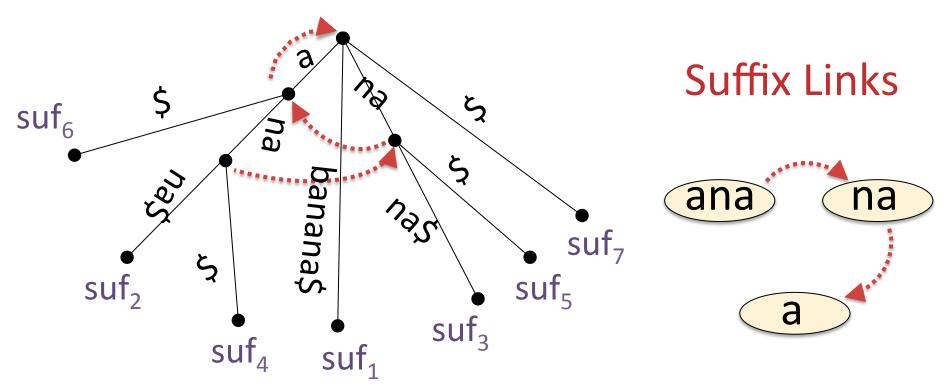


S = banana\$ banana\$ suf₁ anana\$ suf₂ nana\$ suf₃ ana\$ suf₄ na\$ suf₅ a\$ suf₆ \$ suf₇

O(n²) time

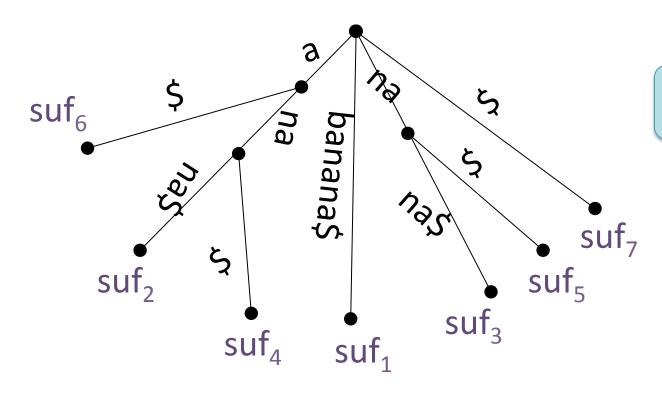


O(n) time



On-line Constructin of Suffix Trees, E. Ukkonen Algorithmica (1995)

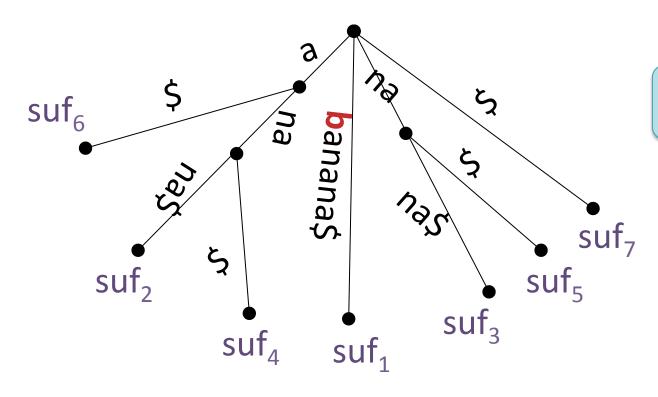




S = banana\$

Search for ban

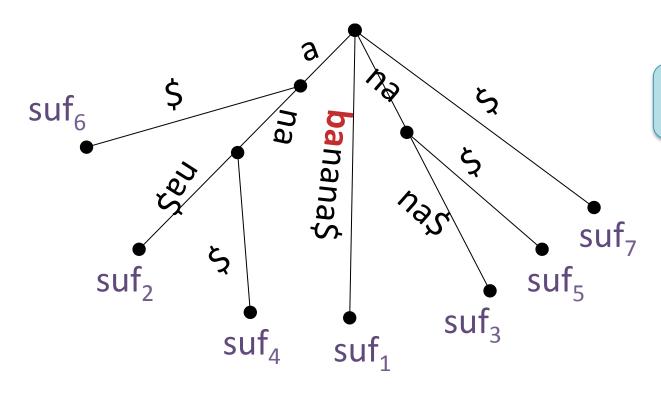




S = banana\$

Search for ban

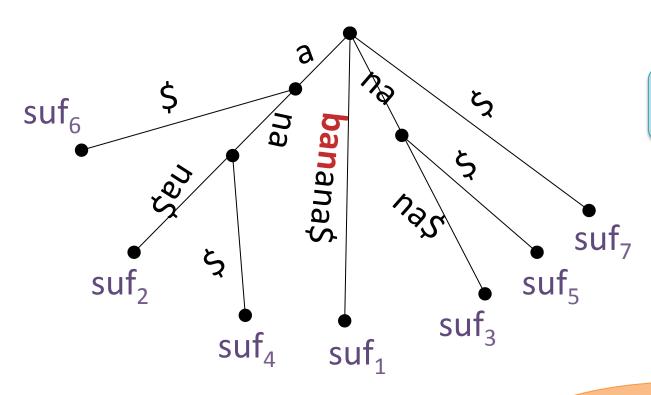




S = banana\$

Search for ban



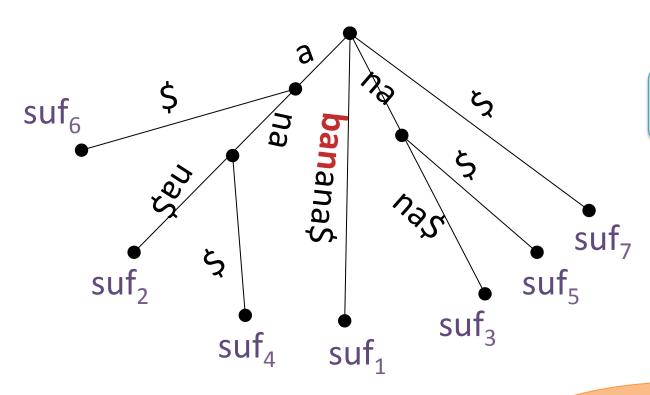


S = banana\$

Search for ban

Found 1 occurrence



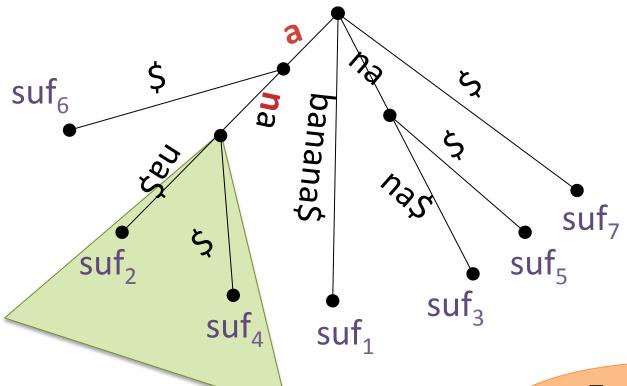


S = banana\$

Search for band

Not found





S = banana\$

Search for an

Found 2 occurrences

Suffix Tree

- ♦ Many applications in computational biology
- ♦Linear time construction algorithms

Linear time solutions to

- Genome alignment
- Finding longest common substring
- All-pairs suffix-prefix matching
- Locating all maximal repetitions
- And many more...

MEMs

Maximal Exact Match (MEM)

Exact match within sequence that cannot be extended left or right without introducing mismatch.



We are interested in MEMs length ≥ k

MEMs

Maximal Exact Match (MEM)

Exact match within sequence that cannot be extended left or right without introducing mismatch.

MEMs are internal nodes in the suffix tree that have left-diverse descendants.

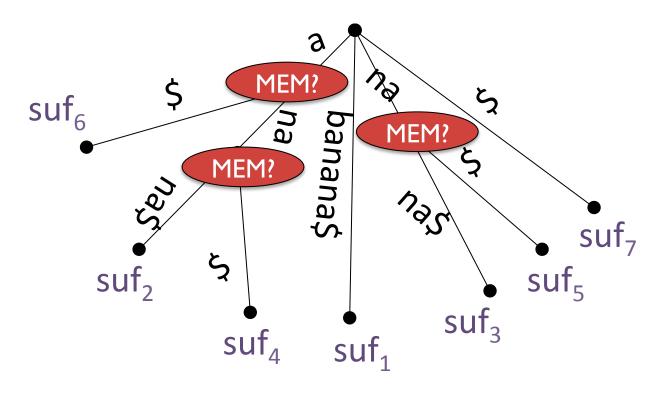
(have descendant leaves that represent suffixes with different characters preceding them)

♦Linear-time suffix tree traversal to locate MEMs.



MEMs in Suffix Tree

Possible MEMs: a, ana, na



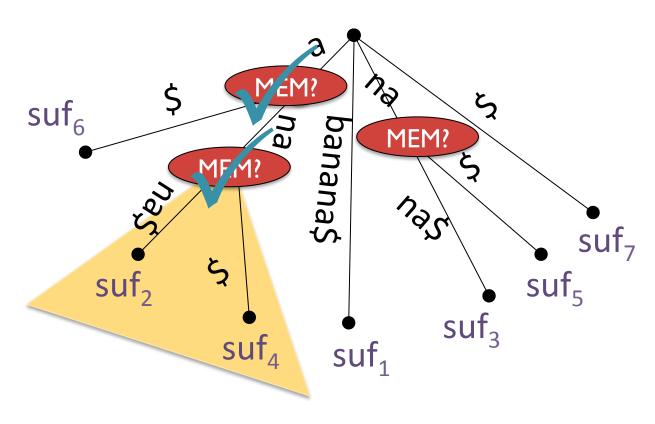
S = banana\$ banana\$ suf₁ anana\$ suf₂ nana\$ suf₃ ana\$ suf₄ na\$ suf₅ a\$ suf₆

MEMs are internal nodes in suffix tree with left-diverse descendants



MEMs in Suffix Tree

MEMs: a, ana



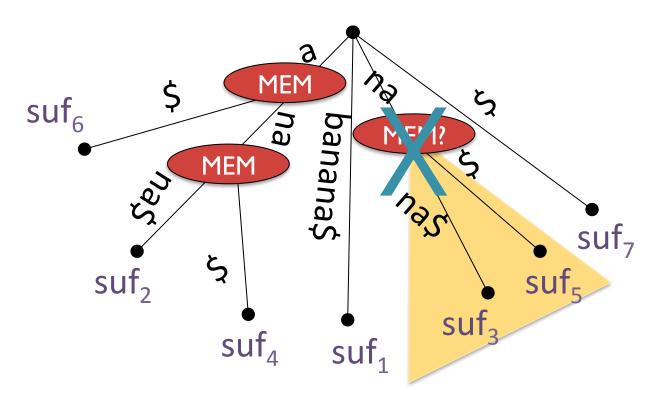
S = banana\$
banana\$ suf₂
nana\$ suf₄

MEMs are internal nodes in suffix tree with left-diverse descendants



MEMs in Suffix Tree

MEMs: a, ana

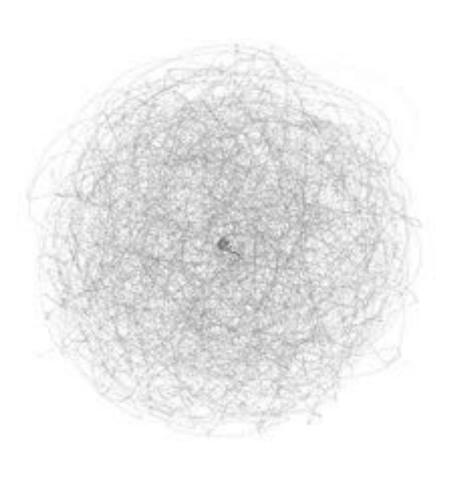


S = banana\$
anana\$ suf₃
ana\$ suf₅

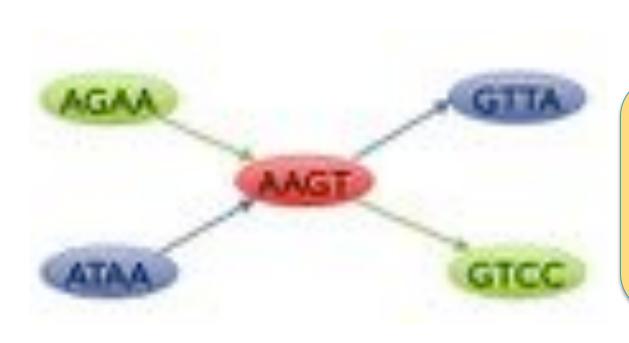
MEMs are internal nodes in suffix tree with left-diverse descendants

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Compresssed de Bruijn graph



Types of nodes:

- i. repeatNodes
- ii. uniqueNodes

Input: AGAAGTCC\$ATAAGTTA

splitMEM

Nodes in compressed de Bruijn graph classified as

- i. repeatNodes
- ii. uniqueNodes

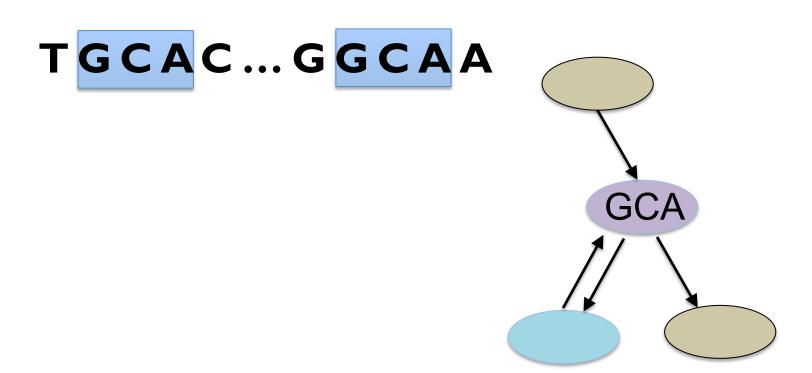
Algorithm:

- I Construct set of repeatNodes
- 2 Sort start positions of repeatNodes
- 3 Create edges and uniqueNodes to link noncontiguous repeatNodes

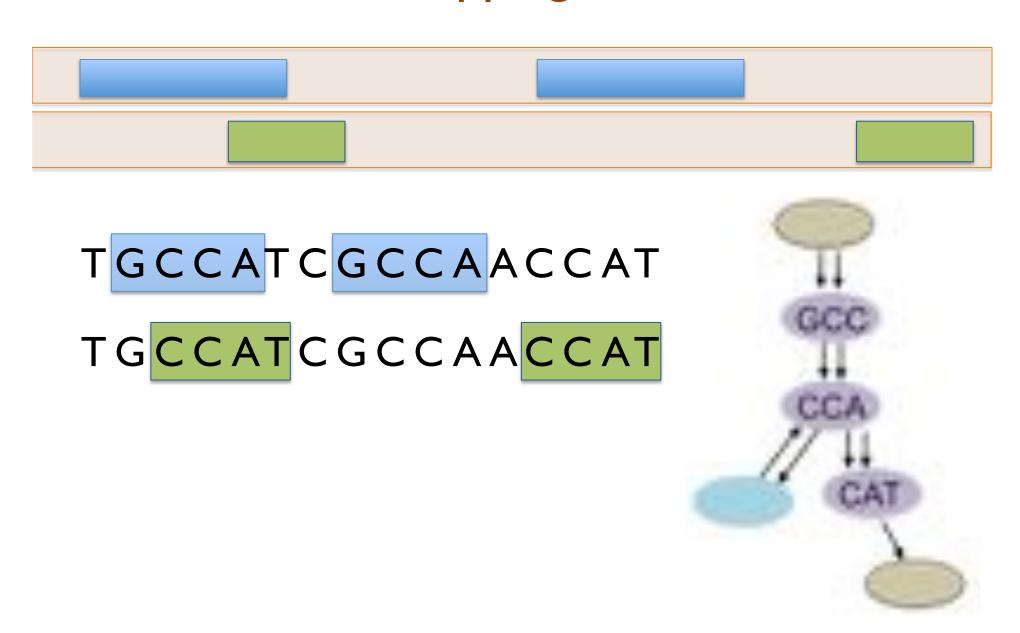
repeatNodes

- I Construct set of repeatNodes
 - I. Build suffix tree of genome
 - 2. Mark internal nodes that are MEMs, length $\geq k$
 - 3. Preprocess suffix tree for LMA queries
 - Compute repeatNodes in compressed de Bruijn graph by decomposing MEMs and extracting overlapping components, length ≥ k

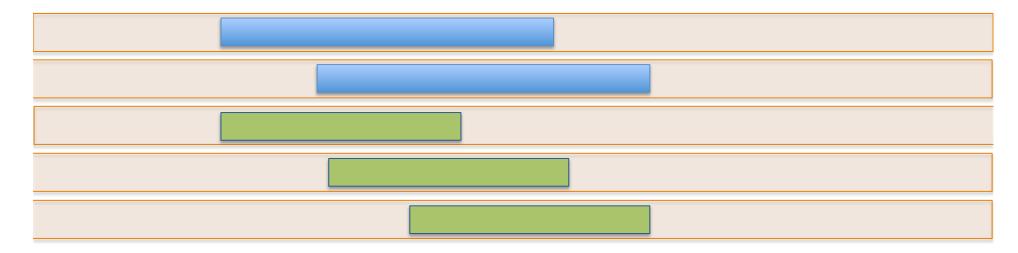
1 MEM occurs twice

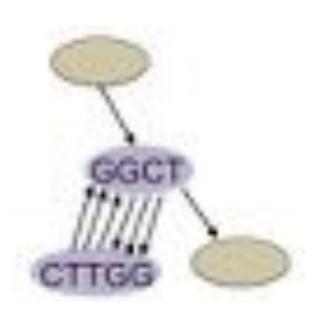


Overlapping MEMs



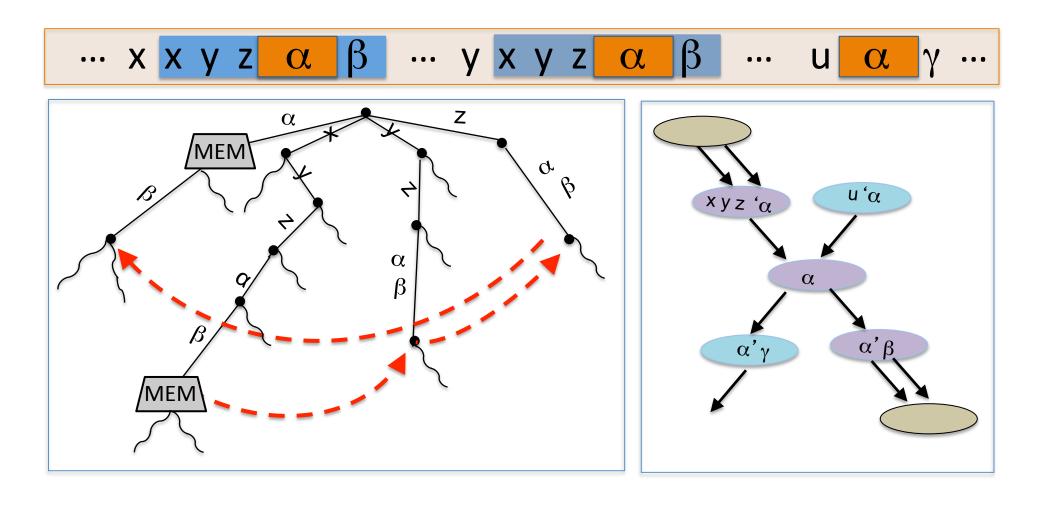
Tandem Repeat

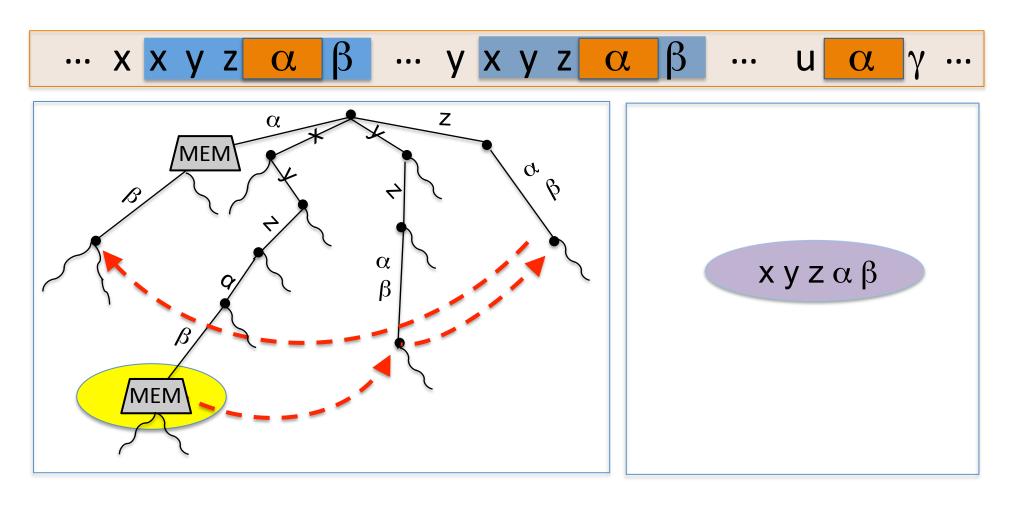




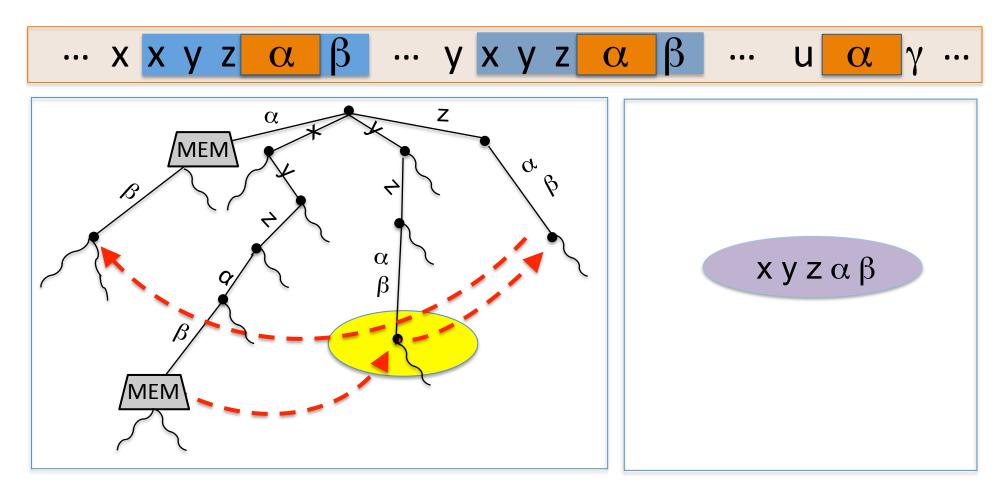
repeatNodes

- I Construct set of repeatNodes
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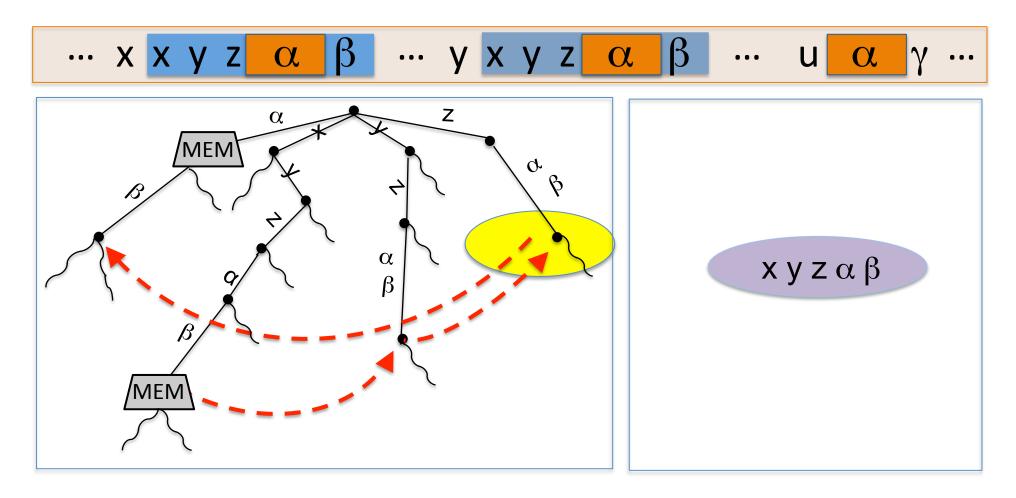


Find MEM in suffix tree.



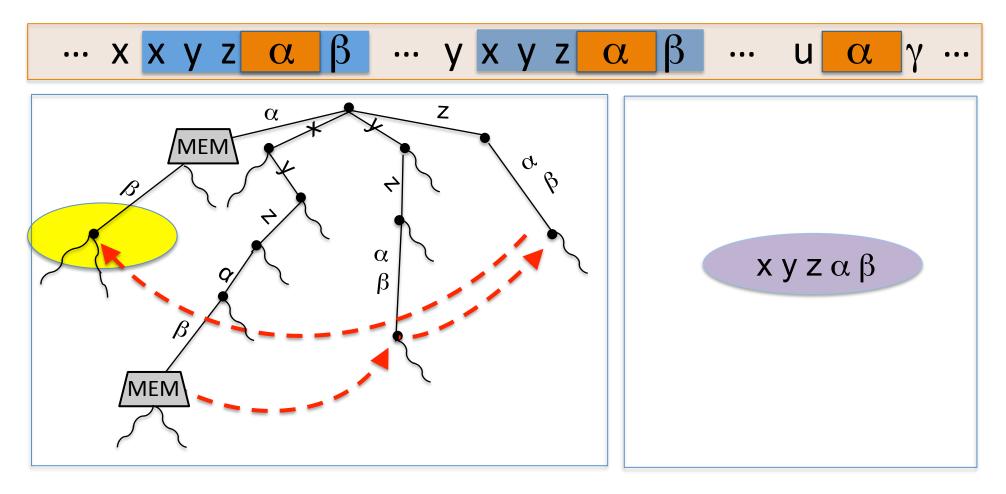
Traverse suffix link.

Look for MEM as ancestor.



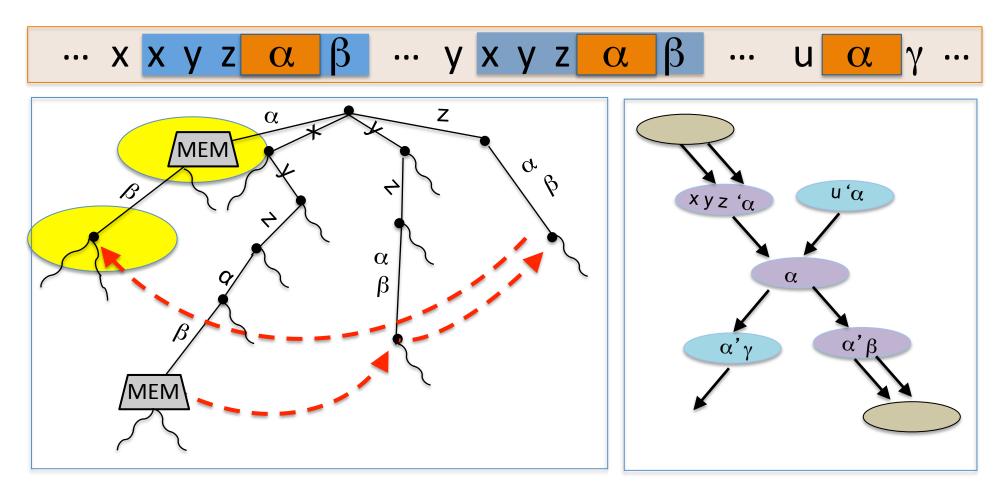
Traverse suffix link.

Look for MEM as ancestor.



Traverse suffix link.

Look for MEM as ancestor.



Found MEM as ancestor. Decompose.

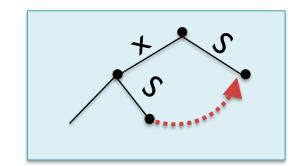
Remove embedded MEM (suffix links). Find next embedded MEM.

Suffix Skips

 \Rightarrow Reduce $O(n^2)$ time to $O(n \log n)$ time

Suffix link: quickly navigate to distant part of tree

- Pointer from internal node labeled xS to node S
- Trim 1 character in O(1) time
- Trim c characters in O(c) time

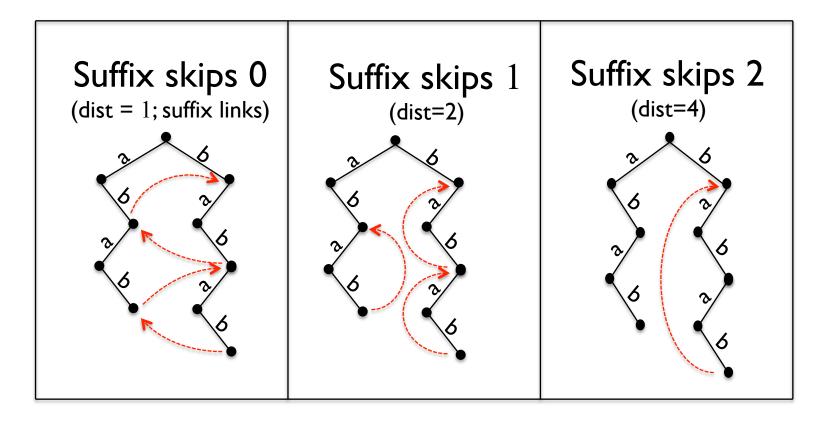


Suffix skip:

Trim c characters in O(log c) time

Suffix Skips

Genome: babab



Additional Preprocessing: pointer jumping to rapidly add additional links

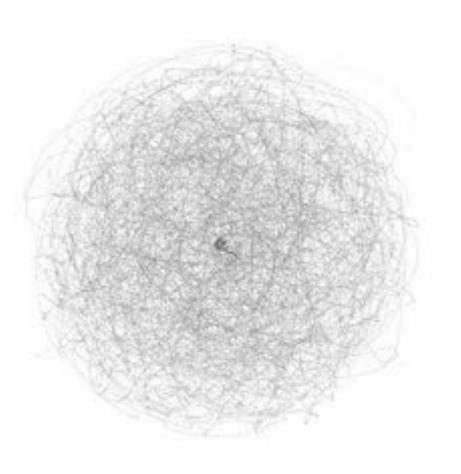
splitMEM

- splitMEM software
 - 0 C++
 - open source http://splitmem.sourceforge.net
- Input modes:
 - o single genome: fasta file
 - o pan-genome: multi-fasta file
- Multi k-mer

construct several compressed de Bruijn graphs without rebuilding suffix tree

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B. Anthracis and E. coli

Examine graph properties:

- Number nodes, edges, avg. degree
- Node length distribution
- Genome sharing among nodes
- Distribution of node distances to core genome

Other properties that can be studied:

- Girth, Diameter, Modularity, Network Motifs, etc.
- Functional enrichment of highly conserved or genome specific genes.

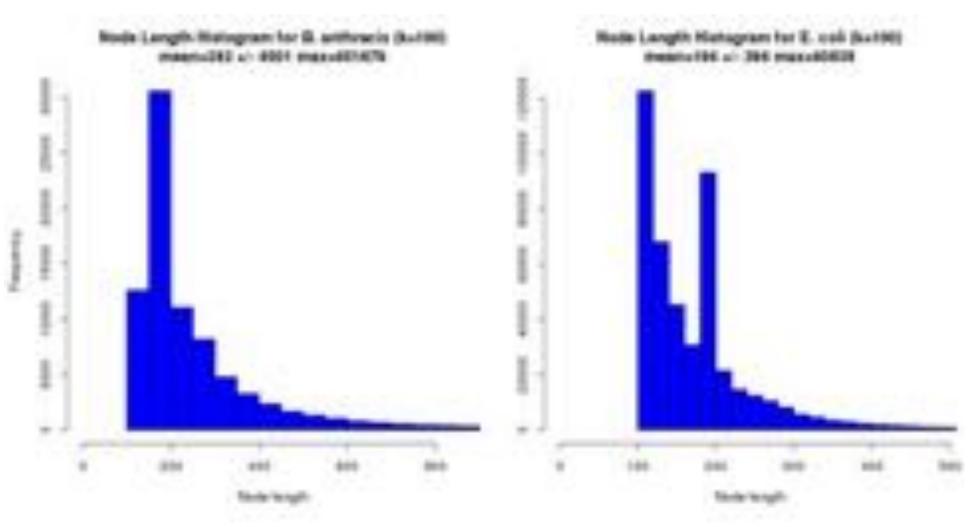
Sean	Sw.	Accrease
B. ambracis A0048 aud/0543	5076 K/A	CP001509
B. anthracio A16R au440003	5179 KB	CP001974
B: anthracis A14 sid40000	\$129.636	CP901970
B. andowin Amer 0581 vid10784	5178 KB	AERCTEM
E. ambracio Attun sici209	5179 KB	ABBINKT9
B. anthracia CDC 486 skE11129	2181 KW	C99001215
R. ambracts Clinic/MAP .	5147 KB	CP00r1av
B. androcis 189001 aid-19901	5170 KB	CF002000
B. arthrecis or Some autobris.	5190 KB	AMM17225
E. com DLZT fits E.Z.Salt ed and CZTT.	2019 K/S	PMINDWA
E. unit 042 sub10ted7		FNSS4Tee
E. anti-536 satte/35	4395 436	CP008247
E cod 55989 aug/3413	51007 K/B	C13939145
E. colli ABU KHITZ widDITZS	50G KX	CP900470
E-onli APEC OL sidrattre	2004 K/K	C7900MM
E 106 APEC OTS wat 6098	4055 KB	C7004009
E. coli 88.21 DE3 (adD0713	4514 KH	CP001509
E. coli ME.21 DED 46429965	450 s.KH	AMNORE

Graphs of main chromosomes

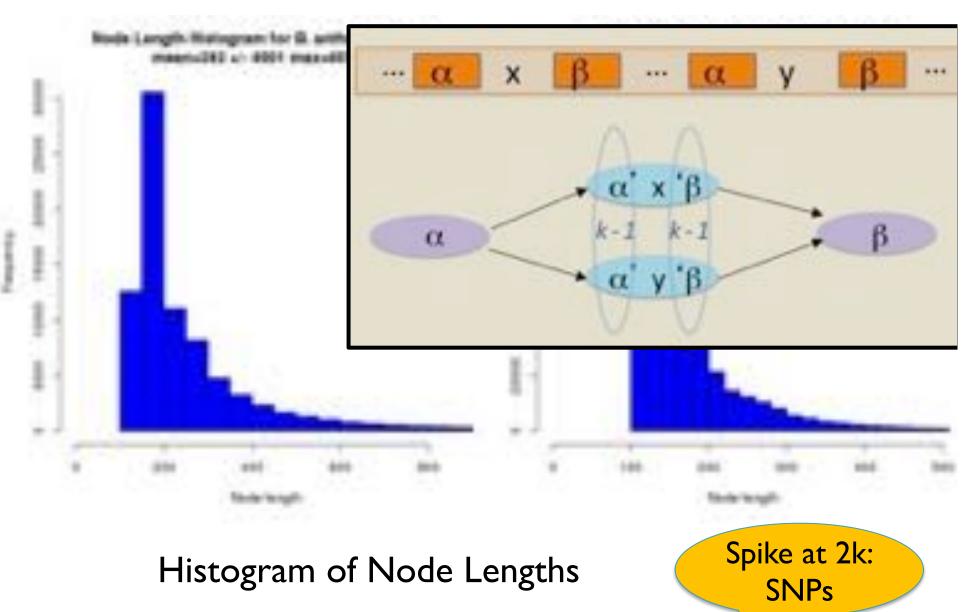
- 9 strains of Bacillus anthracis
- Selection of 9 strains of Escherichia coli

Species	K	Nodes	Edges	Avg. Degree
B. anthracis	25	103926	138468	1.33
B. anthracis	100	41343	54954	1.32
B. anthracis	1000	6627	8659	1.30
E. coli	25	494783	662081	1.33
E. coli	100	230996	308256	1.33
E. coli	1000	11900	15695	1.31

- B. Anthracis and E. coli
- Examine graph properties
- Node length distribution
- Genome sharing among nodes
- Distribution of node distances to core genome



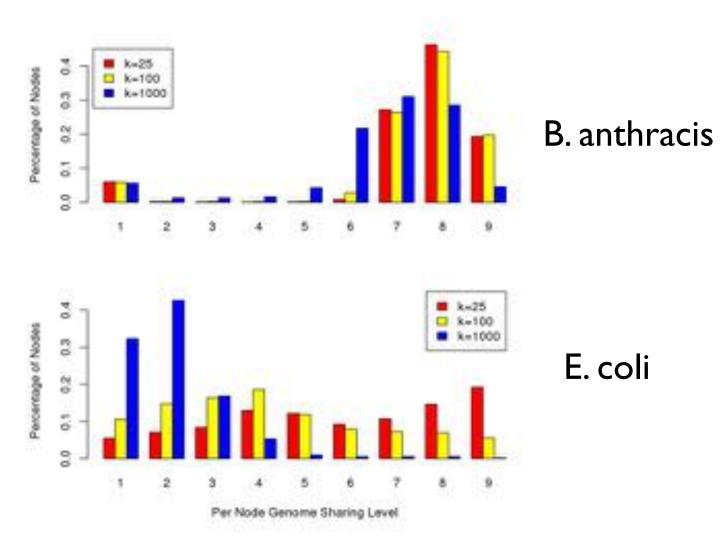
Histogram of Node Lengths



B. Anthracis and E. coli

Examine graph properties

- Node length distribution
- Genome sharing among nodes
- Distribution of node distances to core genome



Fraction of nodes with each level of genome sharing

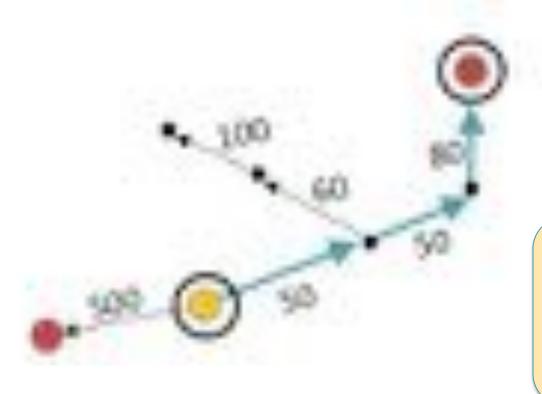
B. Anthracis and E. coli

Examine graph properties

- Node length distribution
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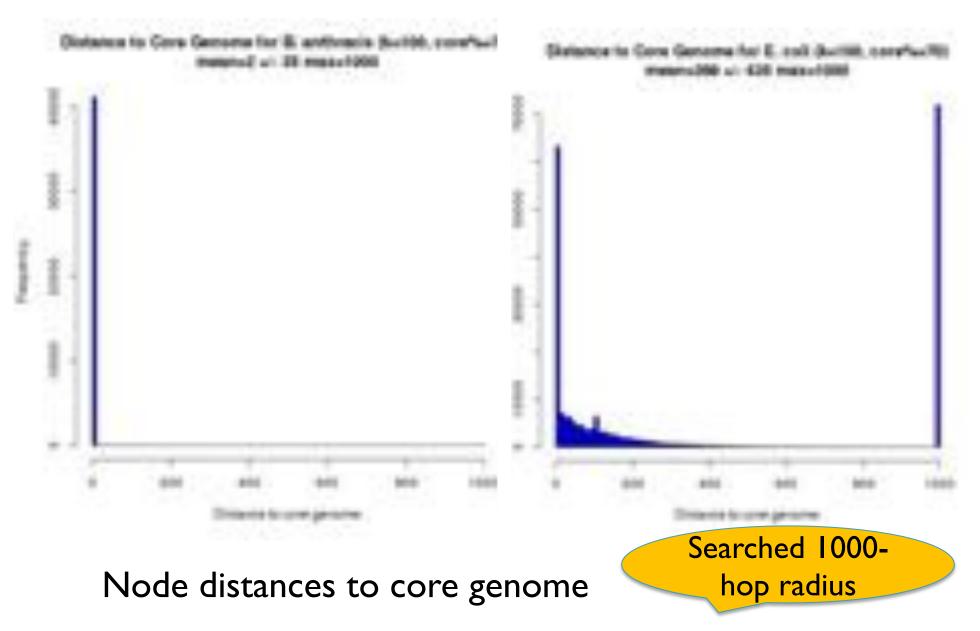
Graph encodes sequence context of segments.

Core genome: subsequences that occur in at least 70% of underlying genomes.



Branch and Bound Search

Nodes can be further in terms of hops while closer by base pairs.



Summary

- Identify pan-genome relationships graphically.
- Topological relationship between suffix tree and compressed de Bruijn graph.
- Direct construction of compressed de Bruijn graph for single or pan-genome.
- Introduce suffix skips.
- Explore pan-genome graphs of B. anthracis, E. coli.

SplitMEM: Graphical pan-genome analysis with suffix skips.

Marcus, S, Lee, H, Schatz, MC (2014) BioRxiv

http://biorxiv.org/content/early/2014/04/06/003954

Future work

Improve splitMEM software:

- Reduce space using compressed full-text index instead of suffix tree
- Approximate indexing of strains to form a pan-genome graph
- Alignment of reads to pan-genome

Biological applications:

- Functional enrichment of core-genome and genome specific segments
- Expand study to larger collection of microbes and larger genomes

Acknowledgments

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Schatz Lab

IT department

Todd Heywood





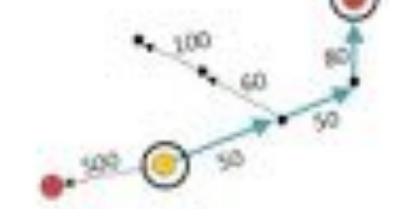
Thank You!

Branch and bound search (like Dijkstra's shortest path algorithm) to compute bp distance from each non-core node to core genome:

Traverse all distinct paths from source until

o a core node is reached

OR o current node was visited by a shorter path



Bounded search

once a core node is found, its distance bounds maximum search distance along other paths